

# COUNTING POINTS ON HYPERELLIPTIC CURVES OVER FINITE FIELDS (1-1)

FCAS  $\hookrightarrow \rightsquigarrow$

## Sec 3.1

$K = \text{field}$      $\text{char } K \neq 2$ .    (eg.  $K = \mathbb{F}_q$  or  $\mathbb{Q}$ )

$g \geq 1$

$C/K = \text{smooth alg. curve assoc. to.}$   
affine eq.  $y^2 = f(x)$

$f \in K[x]$  squarefree,  $\deg f = 2g+1$   
or  $2g+2$

$$f(x) = f_0 + f_1 x + \dots + f_{2g+2} x^{2g+2}$$

$f_j \in K$

(allow  $f_{2g+2} = 0$ )

"Hyperell. curve of genus  $g$  over  $K$ "

Suppose  $L/K = \text{finite extn.}$

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Q: describe  $C(L)$ .

Morphism  $x: C \rightarrow \mathbb{P}^1$

Induces  $x: C(L) \rightarrow \mathbb{P}^1(L) = L \cup \{\infty\}$   
(map of sets)

Q: given  $\alpha \in \mathbb{P}^1(L)$ , what are  $P \in C(L)$   
that map to  $\alpha$ ?

i) If  $\alpha \in L$ , just solve  $y^2 = f(x)$  in  $L$ .  
 $\Rightarrow 0, 1, 2$  points.

ii)  $\alpha = \infty$ ? Need to switch models.

$$u = 1/x \quad (\text{on } \mathbb{P}^1)$$

$$x = \infty \iff u = 0.$$

$$v = y/x^{g+1}.$$

$$y^2 = f_0 + f_1 x + \dots + f_{2g+2} x^{2g+2} \quad \underline{L-3}$$

$$v^2 = f_0 u^{2g+2} + f_1 u^{2g+1} + \dots + f_{2g+2}$$

Subst.  $u \neq 0$ .

want to solve  $v^2 = f_{2g+2} =: "f(\infty)"$

Again: 0, 1, 2 solutions.

### SUMMARY (LEM 3.16)

For each  $\alpha \in L \cup \{\infty\}$ ,

# points  $P \in C(L)$  such that  $x(P) = \alpha$

is 0, 1, 2 depending on

# square roots of  $f(\alpha)$  in  $L$ .

SEC 3.2 | Now assume  $K = \mathbb{F}_q$  (1-4)

Then  $\#C(\mathbb{F}_q) < \infty$ .

Goal of course: given  $f \in \mathbb{F}_q[x]$ ,  
Compute  $\#C(\mathbb{F}_{q^r})$   $r=1, 2, \dots$   
"quickly".

Enumeration method: try every  $\alpha \in \mathbb{P}^1(\mathbb{F}_{q^r})$ .

Very slow: Complexity at least  $q^r$ .  
(exp. in  $r$ )

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Naive bound:  $\#C(\mathbb{F}_{q^r}) \leq 2(q^r + 1)$

Hasse-Weil bound:

$$\underbrace{|\#C(\mathbb{F}_{q^r}) - (q^r + 1)|}_{\text{main term}} \leq \underbrace{2g \cdot q^{r/2}}_{\text{error term.}} \quad (r \geq 1)$$

Sec 3.3  $C/K$   $y^2 = f(x)$ .  $(-5)$   $K = \mathbb{F}_q$

Let  $N_r = \#C(\mathbb{F}_{q^r})$ .  $r \geq 1$ .

Zeta fn:

$$Z_C(T) = \exp\left(N_1 T + \frac{N_2}{2} T^2 + \frac{N_3}{3} T^3 + \dots\right) \in \mathbb{Q}[[T]].$$

Weil "Conjectures"

\*  $Z_C(T)$  is a rational fn.

$$Z_C(T) = \frac{L_C(T)}{(1-T)(1-qT)}$$

$$L_C(T) = 1 + a_1 T + a_2 T^2 + \dots + a_{2g} T^{2g}$$

$a_i \in \mathbb{Z}$

\* Functional eqn:  $L_C\left(\frac{1}{qT}\right) = (qT^2)^{-g} L_C(T)$ .

$$\Leftrightarrow a_{2g-i} = q^{g-i} a_i$$

$i = 0, \dots, 2g.$

\* Riemann Hypothesis:

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roots of  $L_c(T)$  have  
Complex absolute value  $q^{-1/2}$ .

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Goal of course: Compute  $L_c(T)$ .  
 $\in \mathbb{Z}[T]$

Using ~~enum. alg.~~ and func. eqns,  
suffices to compute  $N_1, \dots, N_g$ .

Using enum. alg. Complexity  $\approx q^g$ .

exp in  $g$   
exp in  $\log q$ .

Famous unsolved problem:

Given  $f \in \mathbb{F}_q[x]$ , can we compute  
 $L_c(T)$  in polynomial time?

input size:  $g \cdot \log q$

Schoof-Pila  
 $(\log q)^{Cg}$

# Sec 4 | trace formula

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From now on  $q = p$ .

$C/\mathbb{F}_p$  hyp. curve genus  $g \geq 1$ .

$$y^2 = f(x) \quad f \in \mathbb{F}_p[x]$$

Baby trace formula (Prop 4.1.1)

$$\text{let } h = f^{\frac{p-1}{2}} \in \mathbb{F}_p[x]$$

$$\text{Then } \# C(\mathbb{F}_p) \equiv 1 - \sum_{j=1}^g h_{j(p-1)} \pmod{p}$$

( $h_i = \text{coeff of } x^i \text{ in } h(x)$ ).